# Link invariants from finite racks

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#### Abstract

We define ambient isotopy invariants of oriented knots and links using the counting invariants of framed links defined by finite racks. These invariants reduce to the usual quandle counting invariant when the rack in question is a quandle. We are able to further enhance these counting invariants with 2-cocycles from the coloring rack's second rack cohomology satisfying a new degeneracy condition which reduces to the usual case for quandles.

KEYWORDS: Finite racks, rack homology, 2-cocycle invariants, knot and link invariants 2000 MSC: 57M25, 57M27, 17D99

### 1 Introduction

A rack is a generally non-associative algebraic structure whose axioms correspond to blackboard-framed isotopy moves on link diagrams. Racks generalize quandles, an algebraic structure whose axioms correspond to the three Reidemeister moves which combinatorially encode ambient isotopy of knot diagrams.

Given a finite quandle T, the set of quandle homomorphisms from a knot quandle Q(K) to T gives us an easily computed knot invariant, namely its cardinality |Hom(Q(K), T)|. This is the quandle counting invariant, also sometimes called the quandle coloring invariant since each homomorphism  $f: Q(K) \to T$  can be pictured as a "coloring" of the knot diagram assigning to each arc  $x_i$  in a knot diagram the element  $f(x_i) \in T$  such that a quandle coloring condition is satisfied at every crossing. Indeed, Fox 3-coloring is the simplest non-trivial example of a quandle coloring invariant for knots.

If T is a non-quandle rack, the set of colorings of arcs of a link diagram by elements of T satisfying the coloring condition at every crossing is invariant only under blackboard-framed isotopy – type I Reidemeister moves which change the framing of the knot also change the number of colorings. In this paper we will exploit a property of finite coloring racks to define computable invariants of ambient isotopy of knots and links incorporating these framed isotopy coloring invariants. The usual quandle coloring invariants then form a special case of these more general rack coloring invariants.

The paper is organized as follows. In section 2 we review the basics of racks, framed links and virtual links. In section 3 we define finite rack based counting invariants and give some examples. In particular, we show that the polynomial version of the invariant specializes to the simple version and contains more information than the simple version. In section 4 we enhance the rack counting invariants with 2-cocycles in the style of [3]. We provide an example showing that the cocycle-enhanced invariant contains more information than the polynomial rack counting invariant alone. In section 5 we collect questions for future research.

### 2 Basic definitions

In this section we review the basic definitions we will need for the remainder of the paper.

#### 2.1 Racks

We begin with a definition from [6].

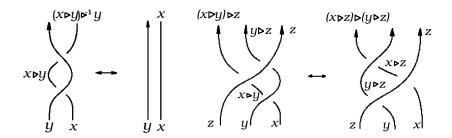
**Definition 1** A rack is a set R with a binary operation  $\triangleright : R \times R \to R$  satisfying

- (i) for all  $x \in R$ , the map  $f_x : R \to R$  defined by  $f_x(y) = y \triangleright x$  is invertible, with inverse  $f_x^{-1}(y)$  denoted  $y \triangleright^{-1} x$ , and
- (ii) for all  $x, y, z \in R$ , we have  $(x \triangleright y) \triangleright z = (x \triangleright z) \triangleright (y \triangleright z)$ .

A rack in which  $x \triangleright x = x$  for all  $x \in R$  is a *quandle*. The operation  $\triangleright^{-1}$  is the *dual* rack operation – it is also self-distributive, and the two operations are mutually distributive. Note that in [6],  $x \triangleright y$  is denoted  $x^y$  and  $x \triangleright^{-1} y$  is denoted  $x^{\overline{y}}$ .

$$\overline{x}$$
  $y$   $\overline{x \triangleright y}$   $y$ 

The rack axioms correspond to Reidemeister moves II and III where we think of rack elements as arcs in an oriented link diagram and  $\triangleright$  means crossing under from right to left when looking in the positive direction of the overcrossing strand. The dual operation  $\triangleright^{-1}$  can then be interpreted as crossing under from left to right.



**Example 1** Perhaps the simplest example of a rack structure on a finite set  $R = \{x_1, \ldots, x_n\}$  is the constant action rack or permutation rack on R associated to a permutation  $\sigma \in S_n$ . Specifically, set

$$x_i \triangleright x_j = x_{\sigma(i)}$$

for all  $i=1,\ldots,n$ ; then the action of  $y\in R$  on R remains constant as y varies. It is easy to verify that this definition gives us a rack structure, since  $x_i \triangleright^{-1} x_j = x_{\sigma^{-1}(i)}$  and we have

$$(x_i \triangleright x_i) \triangleright x_k = x_{\sigma^2(i)} = (x_i \triangleright x_k) \triangleright (x_i \triangleright x_k).$$

If a constant action rack is a quandle, then we have  $x \triangleright x = x$  and consequentially  $x \triangleright y = x$  for all  $x, y \in R$ ; such a quandle is called *trivial*. There is one trivial quandle for each cardinality n, denoted  $T_n$ . We will denote the constant action rack associated to  $\sigma \in S_n$  by  $T_{\sigma}$ .

**Example 2** A simple example of a nontrivial rack structure from [6] is the (t, s)-rack structure: let  $\ddot{\Lambda}$  be the ring  $\mathbb{Z}[t, t^{-1}, s]$  modulo the ideal generated by  $s^2 - (1 - t)s$ . Then any  $\ddot{\Lambda}$ -module M is a rack under the operation

$$x \triangleright y = tx + sy$$
.

For instance, we can take  $M = \mathbb{Z}_n$  and choose  $t, s \in M$  such that gcd(n, t) = 1 and  $s^2 = (1 - t)s$ , e.g.  $M = \mathbb{Z}_8$  with t = 3 and s = 2. If s = 1 - t then M is a quandle, known as an Alexander quandle.

One useful way to describe a rack operation  $\triangleright$  on a finite set  $\{x_1, \ldots, x_n\}$  is to encode its operation table as a matrix M whose entry in row i column j is k where  $x_k = x_i \triangleright x_j$ . Thus, the constant action rack on  $R = \{x_1, x_2, x_3\}$  defined by  $\sigma = (123)$  has matrix

$$M_{(123)} = \left[ \begin{array}{rrr} 2 & 2 & 2 \\ 3 & 3 & 3 \\ 1 & 1 & 1 \end{array} \right]$$

and the (t, s)-rack  $M = \mathbb{Z}_8$  with t = 3 and s = 2 has rack matrix

$$M_{(\mathbb{Z}_8,3,2)} = \begin{bmatrix} 5 & 7 & 1 & 3 & 5 & 7 & 1 & 3 \\ 8 & 2 & 4 & 6 & 8 & 2 & 4 & 6 \\ 3 & 5 & 7 & 1 & 3 & 5 & 7 & 1 \\ 6 & 8 & 2 & 4 & 6 & 8 & 2 & 4 \\ 1 & 3 & 5 & 7 & 1 & 3 & 5 & 7 \\ 4 & 6 & 8 & 2 & 4 & 6 & 8 & 2 \\ 7 & 1 & 3 & 5 & 7 & 1 & 3 & 5 \\ 2 & 4 & 6 & 8 & 2 & 4 & 6 & 8 \end{bmatrix}.$$

Rack axiom (i) requires the columns of a rack matrix to be permutations. See [8] for more.

#### 2.2 Framed links

Recall that a framed link is a link L with a choice of framing curve  $F_i$  for every component  $C_i$  of L, i.e.  $F_i$  is a longitude of a regular neighborhood of  $C_i$ . Framing curves are determined up to isotopy by their linking numbers with  $C_i$ . In terms of diagrams, we can bestow a canonical framing on every component of a link via the blackboard framing, i.e. drawing a framing curve for each  $C_i$  parallel to  $C_i$ . This gives a framing with linking number given by the self-writhe  $sw(C_i) = \sum_{x \in S_i} sign(x)$  where  $S_i$  is the set of crossings where  $C_i$  crosses itself, sign() = 1 and sign() = -1.

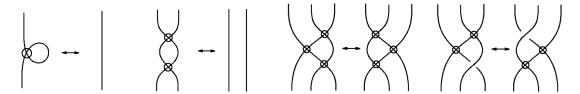
Combinatorially, blackboard-framed links can be regarded as equivalence classes of link diagrams under the equivalence relation generated by Reidemeister moves II and III together with a doubled type I move which preserves the framing of each component; see [14, 6].

### 2.3 Virtual links

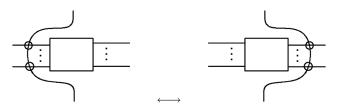
Virtual knot theory is a combinatorial generalization of ordinary classical knot theory; geometrically, a virtual link is an ordinary link in which the ambient space is not  $S^3$  or  $\mathbb{R}^3$  but  $\Sigma \times [0,1]$  for some compact surface  $\Sigma$ , considered up to stabilization (see [10, 4]). More formally, we have:

**Definition 2** A virtual link is an equivalence class of link diagrams with an extra crossing type known as a virtual crossing,  $\boxtimes$ , under the equivalence relation determined by the usual Reidemeister

moves together with the four virtual moves



We can summarize the rules for virtual moves with the *detour move*, which says that any strand with only virtual crossings can be replaced by any other strand with the same endpoints with only virtual crossings. That is, a strand with only virtual crossings can move past any virtual tangle.



Virtual crossings have no intrinsic over- or under-sense, as they are artifacts of drawing non-planar link diagrams on planar paper. Classical links are then virtual links whose underlying surface  $\Sigma$  is  $S^2$ . Replacing the classical Reidemeister I move with the doubled version yields framed virtual links. For the remainder of this paper, we will use "link diagram" to mean "oriented blackboard-framed virtual or classical link diagram." See [10] for more.

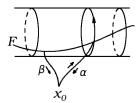
#### 2.4 The fundamental rack(s) of a link

Associated to a framed link L is a rack known as the fundamental rack of L, which we will denote by FR(L) ([6]).

Geometrically, elements of FR(L) are homotopy classes of paths in the link complement  $X = S^3 \setminus (L \times \operatorname{Int}(B^2))$  from the framing curves  $\cup F_i \subset \partial(X)$  to a fixed base point  $x_0$  where the terminal point is fixed but the initial point is permitted to wander along the framing curve  $F_i$  during the homotopy. Any such path  $\alpha : [0,1] \to X$  has an associated element  $\pi(\alpha)$  of the fundamental group  $\pi_1(X,x_0)$  defined by traveling backwards along  $\alpha$ , then going around the canonical meridian in  $\partial(X)$  intersecting  $\alpha(0)$ , then going back along  $\alpha$ . The rack operation is then

$$[\beta] \triangleright [\alpha] = [\beta * \pi(\alpha)]$$

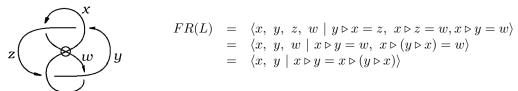
where \* is concatenation of paths.



Combinatorially, given a diagram of L, the fundamental rack of L consists of equivalence classes of rack words in generators corresponding to arcs in the diagram of L under the equivalence relation generated by the rack axioms together with the relations imposed at each crossing. If L is a virtual link, we simply ignore the virtual crossings.

<sup>&</sup>lt;sup>1</sup>Does the notation "FR" stand for "fundamental rack" or "Fenn and Rourke"? Perhaps both!

**Example 3** The pictured blackboard-framed virtual link has fundamental rack with generators x, y and relation  $x \triangleright y = x \triangleright (y \triangleright x)$ :



For each framing of a given link, we have a fundamental rack, generally distinct from the racks of the other framings. All of these racks have a common quotient quandle obtained by setting  $a \triangleright a = a$  for all elements  $a \in FR(L)$ , which is the *knot quandle* Q(L) of the unframed link L. Elements of the knot quandle may be interpreted geometrically as homotopy classes of paths where the initial point is permitted to wander not just along the framing curve but along all of  $\partial(X)$ . See [6] for more.

## 3 Racks and counting invariants

Let L be an unframed link with an ordering on the components. If L has n components, then the framings of L may be indexed by n-tuples  $\mathbf{w} \in \mathbb{Z}^n$ , each with its own a priori distinct fundamental rack. At the most basic level, then, there are infinitely many rack counting invariants for a given link with respect to any choice of finite target rack T. However, we can make a useful observation which enables us to get computable ambient isotopy invariants from the  $\mathbb{Z}^n$ -set of racks of L.

**Definition 3** Let T be a finite rack. For any  $x \in T$ , let  $x^{\triangleright n}$  for  $n \in \mathbb{Z}_+$  be defined recursively by

$$x^{\triangleright 1} \equiv x \triangleright x$$
 and  $x^{\triangleright (k+1)} \equiv x^{\triangleright k} \triangleright x^{\triangleright k}$ .

For each element  $x \in T$ , the *rack exponent* of  $x \in T$ , denoted  $\rho(x)$ , is the minimal natural number  $n \in \mathbb{Z}_+$  such that  $x^{\triangleright n} = x$ . The *rack rank* of T, denoted N(T) or just N if T is understood, is the least common multiple of the rack exponents of the elements of T,

$$N(T) = \operatorname{lcm}\{\rho(x) \mid x \in T\}.$$

To see that  $\rho(x)$  is well defined for all  $x \in T$ , we first need a lemma.

**Lemma 1** Let T be any rack and  $x, y \in T$ . Then  $y \triangleright (x \triangleright x) = y \triangleright x$ .

Proof.

$$y \triangleright (x \triangleright x) = [(y \triangleright^{-1} x) \triangleright x] \triangleright (x \triangleright x) = [(y \triangleright^{-1} x) \triangleright x] \triangleright x = y \triangleright x.$$

**Remark 4** Two elements  $x, y \in T$  are operator equivalent if  $z \triangleright x = z \triangleright y$  for all  $z \in T$ . If T is a finite rack, then two elements are operator equivalent iff their columns in the matrix of T are identical. Lemma 1 says that the  $\triangleright$ -powers of  $x \in T$  are all operator equivalent. Indeed, the set of operator equivalence classes of a rack forms a quandle under the natural operation  $[x] \triangleright [y] = [x \triangleright y]$ .

**Corollary 2** *Let* T *be a rack. If*  $x \triangleright x = y \triangleright y$ , *then* x = y.

**Proof.** Suppose  $x \triangleright x = y \triangleright y = z$ . We have  $x \triangleright x = x \triangleright (x \triangleright x) = x \triangleright z$  and  $y \triangleright y = y \triangleright (y \triangleright y) = y \triangleright z$ . Then  $x \triangleright x = y \triangleright y$  implies  $x \triangleright z = y \triangleright z$  and rack axiom (i) implies x = y.

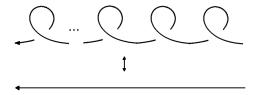
In terms of rack matrices, corollary 2 says that like the columns of a rack matrix, the diagonal of a rack matrix must be a permutation. It then easily follows that  $\rho(x) < |T|$  for any  $x \in T$  where T is a finite rack – indeed, N(T) is just the exponent of the permutation along the diagonal. This fact also follows from proposition 7.3 in [12].

We will also need the following standard result (see [6] or [14] for example):

**Theorem 3** If D and D' are ambient isotopic link diagrams, we can modify D' to obtain a diagram D'' which is framed isotopic to D by selecting an arc on each component of D' and adding positive or negative kinks until the framings match.

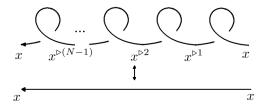
The proof of theorem 3 involves taking any Reidemeister move sequence starting with D and ending with D' and replacing every type I move with a double I move to adjust the framed isotopy class; at the end, we can then slide the extra crossings along the component until they arrive at the chosen arc. Note that this argument applies to virtual links as well as classical links, since we can slide a classical kink past a virtual crossing using a detour move. Note also that without loss of generality we can assume that all kinks added have positive winding number since we need not preserve the regular isotopy class, only the blackboard-framed isotopy class.

**Definition 4** Let  $N \in \mathbb{N}$ . We say two blackboard-framed oriented link diagrams are N-phone cord equivalent if one may be obtained from other by a finite sequence of Reidemeister II and III moves and the following N-phone cord move, where N is the number of loops:



**Proposition 4** Let T be a finite rack with rack rank N. If two link diagrams D and D' are N-phone cord isotopic then  $|\operatorname{Hom}(FR(D),T)| = |\operatorname{Hom}(FR(D'),T)|$ .

**Proof.** From the definition of rack rank, it is easy to see that N-phone cord moves induce a bijection on the set of colorings as illustrated.



For two *n*-tuples  $\mathbf{v}$ ,  $\mathbf{w} \in \mathbb{Z}^n$ , let us write  $\mathbf{v} \equiv \mathbf{w} \mod N$  if for all components  $i = 1, \dots n$  we have  $v_i \equiv w_i \mod N$ .

**Corollary 5** Let T be a finite rack with rack rank N. If two link diagrams D and D' are ambient isotopic and have self-writhe vectors congruent modulo N, then |Hom(FR(D), T)| = |HomFR(D'), T|.

Note that if T is a finite rack with rack rank N and L is a link, the set of self-writhes of each component of L modulo N can be indexed by  $\mathbf{w} \in (\mathbb{Z}_N)^c$  where c is the number of components of L. For ease of notation, when N and c are understood let us denote  $(\mathbb{Z}_N)^c = W$  and a blackboard-framed diagram of D with self-writhe vector  $\mathbf{w} \in W$  by  $(D, \mathbf{w})$ .

We can now define computable unframed knot and link invariants using these cardinalities.

6

**Definition 5** Let T be a finite rack and L a link with c components. The *integer rack counting invariant* of L with respect to T is given by

$$IR(L,T) = \sum_{\mathbf{w} \in W} |\mathrm{Hom}(FR(D,\mathbf{w}),T)|.$$

Note that if T is a quandle, then we have N(T) = 1 and IR(L,T) is the ordinary quandle counting invariant |Hom(Q(L),T)|. Hence the integer rack counting invariant is the natural generalization of the quandle counting invariant to the finite rack case.

**Example 5** If  $T = \{x_1, \ldots, x_n\}$  is a constant action rack defined by an n-cycle, an undercrossing color  $\tau$  becomes  $\sigma(\tau)$  if going right-to-left and  $\sigma^{-1}(\tau)$  if going left-to-right, so pushing a color around the knot yields an ending color of  $\sigma^{writhe(K)}(\tau)$  if our starting color was  $\tau$ . Thus, there is a rack coloring of a framed knot K by T if and only if the writhe of K is zero mod n. Indeed, there are n such colorings for the 0-framing mod n and none for the others, and we have SR(K,T) = n + (n-1)0 = n for any knot K. This generalizes the fact that |Hom(K,T)| = n for T a trivial quandle of cardinality n and K a knot.

**Example 6** Let  $T_{Ex6}$  be the rack with matrix

The integer rack counting invariant with respect to  $T_{Ex6}$  distinguishes the trefoil  $3_1$  from the unknot  $U_1$  with  $IR(U_1, T_{Ex6}) = 10$  and  $IR(3_1, T_{Ex6}) = 22$ , as the reader can verify from the tables of colorings listed in table 1. Here  $N(T_{Ex6}) = 2$ , so we need only consider one diagram each of  $U_1$  and  $3_1$  with odd writhe and one of each with even writhe.

We can enhance the integer rack counting invariant by keeping track of which framings contribute which colorings. For a writhe vector  $\mathbf{w} = (w_1, \dots, w_c) \in W = (\mathbb{Z}_N)^c$  let us denote the product  $\prod_{k=1}^c q_1^{w_1} q_2^{w_2} \dots q_c^{w_c}$  by  $q^{\mathbf{w}}$ . Then we have:

**Definition 6** Let T be a finite rack and L a link with c components and self-writh vector  $\mathbf{w} = (w_1, \dots, w_c) \in W$ . The polynomial rack counting invariant of L with respect to T is given by

$$PR(L,T) = \sum_{\mathbf{w} \in W} |\mathrm{Hom}(FR(L,\mathbf{w}),T)| q^{\mathbf{w}}.$$

The polynomial version of the rack counting invariant holds more information than the simple version, as the next example shows.

**Example 7** The constant action rack T with rack matrix  $\begin{bmatrix} 2 & 2 \\ 1 & 1 \end{bmatrix}$  has rack rank N(T) = 2. The Hopf link H and the two-component unlink  $U_2$  both have integer rack counting invariant value 4 with respect to T, but their polynomial rack counting invariants are distinct, with  $PR(H,T) = 4q_1q_2$  and  $PR(U_2,T) = 4$  as the reader can easily verify from table 2.

Indeed, generalizing the preceding example we have

D	Colorings by $T_{Ex6}$
y $y$ $y$ $y$ $y$ $y$ $y$ $y$ $y$ $y$	$ \begin{bmatrix} x & 1 & 1 & 1 & 2 & 2 & 2 & 3 & 3 & 3 \\ y & 1 & 2 & 3 & 1 & 2 & 3 & 1 & 2 & 3 \\ z & 1 & 3 & 2 & 3 & 2 & 1 & 2 & 1 & 3 \end{bmatrix} $
y $w$	$ \begin{array}{ c cccccccccccccccccccccccccccccccccc$
$\int_{X}$	

Table 1: Rack colorings of  $\mathbf{3}_1$  and  $U_1$  by  $T_{Ex6}$ 

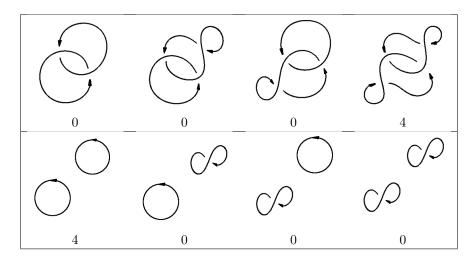


Table 2: Numbers of colorings of H and  $U_2$  by  $T_{(12)}$ .

**Proposition 6** Let L be a two-component classical link and  $T = T_{\sigma}$  a constant action rack with  $\sigma \in S_N$  an N-cycle. Then the polynomial rack counting invariant has the form

$$PR(L,T) = N^2 q_1^l q_2^l$$

where l is the negative of the linking number  $lk(L_1, L_2)$  of L mod N.

**Proof.** Traveling around a component, to get a valid coloring the end color must match the initial color, so we must go through N crossings (counted algebraicially). Since  $lk(L_1, L_2)$  of these are multi-component crossings which do not contribute to the self-writhe, we must have

$$l + lk(L_1, L_2) = N.$$

The same holds for both components if L is classical. There are N choices of starting color for each component and every pair produces exactly one coloring, so there are  $N^2$  total colorings.

Corollary 7 If L is a two-component classical link and  $T = T_{\sigma}$  a constant action rack with  $\sigma \in S_N$  an N-cycle such that the exponents of of  $q_1$  and  $q_2$  differ in any term of PR(L,T), then L is non-classical.

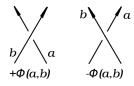
# 4 Rack cocycle invariants

In this section we generalize the quandle 2-cocycle invariants defined in [3] to the finite rack case.

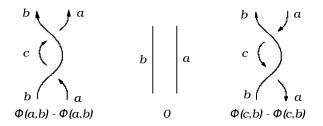
The rack counting invariants described in the last section work by taking cardinalities of sets of homomorphisms which are unchanged by Reidemeister moves. However, a set is more than a cardinality, and we would like to recover as much information from these sets of homomorphisms as possible.

In [3], the idea is to associate a sum in an abelian group A called a *Boltzmann weight* to a quandle-colored knot diagram in such a way that the sum does not change under Reidemeister moves. Then, instead of counting "1" for each homomorphism, we count its Boltzmann weight, transforming the set of colorings into a multiset of these weights. Such multisets are commonly encoded as polynomials by converting the multiset elements to exponents and multiplicities to coefficients of a dummy variable, e.g.  $\{1, 1, 1, 4, 4\}$  becomes  $3t + 2t^4$ .

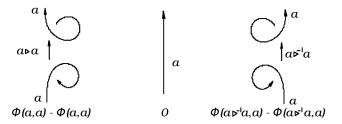
The Boltzmann weights are defined as follows: at every crossing in a rack-colored link diagram, we want to count  $\phi(a, b)$  at a positive crossing or  $-\phi(a, b)$  at a negative crossing where b is the color on the overcrossing strand and a is the color on the inbound understrand for positive crossings and the outbound understrand for negative crossings.



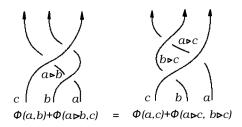
This weighting rule has the advantage that the contributions from the two crossings in a Reidemeister type II move cancel, so the sum is automatically invariant under II moves:



We also note that the weighting rule gives invariance under the doubled type I moves required for blackboard-framed isotopy:



The condition for the sum to be unchanged by Reidemeister III moves is pictured below.



This turns out to be the condition that  $\phi$  is a cocycle in the second rack cohomology  $H^2_R(T;A)$  of the rack T with coefficients in A. Specifically, the A-module spanned by n-tuples of elements of T is the space of rack n-chains  $C_n^R(T;A) = A[T^n]$ ; its dual is the space of rack n-cochains  $C_n^R(T;A) = \text{Hom}(C_n^R(T;A),A)$ . Note that  $C_n^R(T;A)$  has A-generating set  $\{\chi_{\mathbf{x}} \mid \mathbf{x} \in T^n\}$  where

$$\chi_{\mathbf{x}}(\mathbf{y}) = \begin{cases} 1 & \mathbf{x} = \mathbf{y} \\ 0 & \text{otherwise} \end{cases}$$

for  $\mathbf{y} \in T^n$ . Next, we define a coboundary map  $\delta^n : C_R^n(T;A) \to C_R^{n+1}(T;A)$  by

$$(\delta^{n}\phi)(x_{1},\ldots,x_{n+1}) = \sum_{i=1}^{n+1} (-1)^{i-1} \left(\phi(x_{1},\ldots,x_{i-1},x_{i+1},\ldots,x_{n+1})\right) - \phi(x_{1} \triangleright x_{i},\ldots,x_{i-1} \triangleright x_{i},x_{i+1},\ldots,x_{n+1})\right).$$

Then for  $\phi$  to yield a Boltzmann weight, we need  $\phi \in \text{Ker}(\delta^2)$ .

To get invariance under the Reidemeister I move, in the quandle case we require that  $\phi(x,x) = 0$  for all  $x \in T$ . This condition also turns out to have a homological interpretation – the cocycles we want to kill live in a subcomplex called the *degenerate* cochains. In the non-quandle rack case, however, a weaker condition suffices.

**Definition 7** Let T be a finite rack with rack rank N, A an abelian group and  $\phi \in C_R^2(T, F)$ . Say  $\phi$  is N-reduced if we have

$$\sum_{k=1}^{N} \phi(a^{\triangleright k}, a^{\triangleright k}) = 0$$

for all  $a \in T$ .

Now we can define an enhanced version of the polynomial rack counting invariant:

**Definition 8** Let L be an oriented blackboard-framed link, T a finite rack and  $\phi$  an N-reduced rack 2-cocycle. For a rack-colored framed diagram of L,  $(D, \mathbf{w})$ , the Boltzmann weight BW(f) of the coloring  $f \in \text{Hom}(FR(D, \mathbf{w}), T)$  is the sum over the crossings in D of the crossing weights,

$$BW(f) = \sum_{c \text{ crossing}} sign(c)\phi(a,b).$$

Then the  $rack\ cocycle\ invariant\ of\ L$  with respect to T is

$$\Phi_{\phi}(L,T) = \sum_{\mathbf{w} \in W} \left( \sum_{f \in \text{Hom}(FR(D,\mathbf{w}),T)} z^{BW(f)} \right) q^{\mathbf{w}}.$$

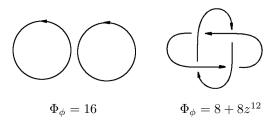
Note that if T is a quandle, then N(T)=1 and  $\phi$  is 1-reduced iff  $\phi(x,x)=0$  for all degenerate cycles  $(x,x)\in C_2^R(T,A)$ ; in this case we also have  $W=\{(0,\ldots,0)\}$ . Indeed, in the quandle case this rack cocycle invariant becomes the usual quandle 2-cocycle invariant from [3].

Specializing z = 1 in  $\Phi_{\phi}(L, T)$  recovers the polynomial rack counting invariant. To see that the rack cocycle invariant is stronger than PR(L, T), consider the following simple example.

#### **Example 8** The rack T with rack matrix

$$M_T = \left[ \begin{array}{cccc} 3 & 1 & 3 & 1 \\ 2 & 4 & 2 & 4 \\ 1 & 3 & 1 & 3 \\ 4 & 2 & 4 & 2 \end{array} \right]$$

has reduced cocycle  $\phi = \chi_{(12)} + \chi_{(14)} + \chi_{(32)} + \chi_{(34)}$  with  $\mathbb{Z}_{13}$  coefficients. Then the (4, 2)-torus link is distinguished from the two-component unlink by  $\Phi_{\phi}$ :

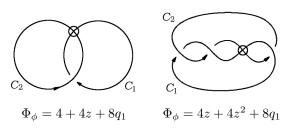


Our final example illustrates a pair of virtual links with equal PR(L,T) values which are distinguished when we include the rack cocycle information.

### **Example 9** Again let T be the rack with rack matrix

$$M_T = \left[ \begin{array}{cccc} 3 & 1 & 3 & 1 \\ 2 & 4 & 2 & 4 \\ 1 & 3 & 1 & 3 \\ 4 & 2 & 4 & 2 \end{array} \right]$$

and  $\phi = \chi_{(1,2)} + \chi_{(1,4)} + \chi_{(3,2)} + \chi_{(3,4)} \in C_R^2(T; \mathbb{Z}_{13})$ . Then  $\Phi_{\phi}$  distinguishes the two pictured virtual links, both of which have  $PR(L,T) = 8 + 8q_1$ . Note that the subscripts on q correspond to the component ordering.



### 5 Questions

In this section we collect a few questions for future research.

Rack and quandle (co)homology has been generalized in various ways including twisted quandle (co)homology in [2], quandle (co)homology with coefficients in quandle modules in [1] and more. How does the rack cocycle invariant change in these cases?

Quandle 3-cocycles have been used to enhance quandle counting invariants of surface knots, i.e. embeddings of compact orientable 2-manifolds in  $S^4$ . How do the rack counting and cocycle invariants extend to the surface knot case?

Other ways of enhancing the quandle counting invariants include using quandle polynomials and exploiting any extra structure the quandle may have (symplectic vector space, *R*-module, etc.); generalizing these ideas to the rack case will be the subject of future papers.

Replacing the arcs in the combinatorial motivation for the rack axioms with semiarcs yields biracks, also known as invertible switches or Yang-Baxter sets (see [7, 5]). The birack analogues of the simple and polynomial rack counting invariants will be the subject of another future paper.

Python code for computing rack counting invariants, reduced rack 2-cocycles with  $\mathbb{Z}_n$  coefficients, and rack cocycle invariants is available for download from www.esotericka.org.

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